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TREE DOMATIC NUMBER IN GRAPHS

Abstract. A dominating set S in a graph G is a tree dominating set of G if the subgraph induced by S is a tree. The tree domatic number of G is the maximum number of pairwise disjoint tree dominating sets in $V(G)$. First, some exact values of and sharp bounds for the tree domatic number are given. Then, we establish a sharp lower bound for the number of edges in a connected graph of given order and given tree domatic number, and we characterize the extremal graphs. Finally, we show that a tree domatic number of a planar graph is at most 4 and give a characterization of planar graphs with the tree domatic number 3.

Keywords: tree domatic number, regular graph, planar graph, Cartesian product.

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1. INTRODUCTION

Let $G = (V, E)$ be a simple graph of order n . The *degree, neighborhood and closed neighborhood* of a vertex v in the graph G are denoted by $d(v)$, $N(v)$ and $N[v] = N(v) \cup \{v\}$, respectively. The *minimum degree and maximum degree* of the graph G are denoted by $\delta(G)$ and $\Delta(G)$, respectively. The graph induced by $S \subseteq V$ is denoted by $\langle S \rangle$. Let P_n , C_n and K_n denote the *path, cycle and complete graph* with n vertices, respectively. Let $K_{n,m}$ denote the *complete bipartite graph*.

A set of vertices D in a graph $G = (V, E)$ is a *dominating set* if every vertex in $V - D$ has at least one neighbor in D . A dominating set D is called a *tree dominating set* if the subgraph induced by D is a tree. In what follows, we assume that all graphs are connected. The minimum number of vertices in a tree dominating set of G is called the *tree domination number* of G , and is denoted by $\gamma_{tr}(G)$. A *tree domatic partition* of G is a partition of the vertex set V into pairwise disjoint tree dominating sets. If such a partition exists, the maximum number of subsets in such a partition is called the *tree domatic number* of G and is denoted by $d_{tr}(G)$, otherwise, define $d_{tr}(G) = 0$.

Zelinka [1] studied the connected domatic number of a graph. Hartnell *et al.* [2] gave the connected domatic number of a planar graph. Chen *et al.* [3] studied tree

dominating set of a graph. We gave some graphs having no tree dominating set. For example, take a complete graph K_{3n} , and partition the vertices of K_{3n} into n sets of three vertices each. To each set of three vertices, add a new vertex and join it to each of these three vertices. Then the resulting 3-connected graph does not have a tree dominating set. If G has no tree dominating set, then we define $d_{tr}(G) = 0$.

A *planar graph* is a graph that can be embedded in the plane so that no two edges intersect geometrically except at a vertex to which both are incident. A planar graph G is called *outerplanar* if the embedding can be chosen so that the boundary of one of the planar regions contains every vertex of G . A graph G_2 is called an *elementary contraction* of G_1 if there is an edge uv of G_1 such that G_2 is obtained from G_1 by deleting the vertices u and v and appending a new vertex, denoted by x , that is adjacent to all the vertices of $G - u - v$ that were originally adjacent to either of u or v . If the graph H is isomorphic to G or is obtainable from G by a finite sequence of elementary contractions, then we say that H is a *contraction* of G . Perhaps a more intuitive way to think of a contraction H of a graph G is to consider a partition of $V(G)$ into subsets each of which induces a connected subgraph of G . Each member of the partition corresponds to a vertex of H , and two vertices of H are adjacent if the union of the corresponding subsets of G induces a connected subgraph of G . In effect each member of the partition has been shrunk to a single vertex and multiple edges have been removed. It is clear that the property of being planar is preserved under contractions.

For an arbitrary graph G , the *vertex connectivity* $\kappa(G)$ is the minimum number of vertices whose removal will disconnect G .

The *Cartesian product* of G and H , denoted by $G \times H$, has the vertex set $V(G \times H) = \{(g, h) | g \in V(G), h \in V(H)\}$ and the edge set $E(G \times H) = \{(g_1, h_1)(g_2, h_2) | \text{either } g_1 = g_2 \text{ and } h_1h_2 \in E(H) \text{ or } g_1g_2 \in E(G) \text{ and } h_1 = h_2\}$. In particular, if G is isomorphic to P_r and H is isomorphic to P_s , then $G \times H$ is called $r \times s$ grid graph, denoted by $G_{r \times s}$.

2. SOME EXACT VALUES OF AND BOUNDS FOR THE TREE DOMATIC NUMBER

First, we give some upper bounds for the tree domatic number in term of minimum degree, tree domination number or vertex connectivity.

Theorem 1. *Let G be a connected graph. Then,*

- (1) *If G is a complete graph, then $d_{tr}(G) = \delta(G) + 1$; otherwise, $d_{tr}(G) \leq \delta(G)$ and the bound is sharp.*
- (2) *If $\gamma_{tr}(G) > 0$, then $d_{tr}(G) \leq \frac{n}{\gamma_{tr}(G)}$ and the bound is sharp.*
- (3) *$d_{tr}(G) \leq \kappa(G)$ and the bound is sharp.*

Proof. (1) Assume $d_{tr}(G) = t$ and (D_1, D_2, \dots, D_t) is a partition of $V(G)$ into t tree dominating sets. If G is a complete graph, then it is obvious that $d_{tr}(G) = n = \delta(G) + 1$. If G is not a complete graph, then let v be a vertex of G such that

$d(v) = \delta(G) < n - 1$. Assume $v \in D_1$. If $|D_1| = 1$, then $D_1 = \{v\}$. Since each vertex in D_i is dominated by v for $i = 2, 3, \dots, t$, it follows that $d(v) = n - 1$, which is a contradiction. Hence, $|D_1| \geq 2$. Since $\langle D_1 \rangle$ is a tree and v is dominated by at least one vertex of D_i for $i = 2, 3, \dots, t$, it follows that $d(v) - 1 \geq t - 1$. That is $t \leq d(v)$. Hence, $d_{tr}(G) \leq \delta(G)$.

(2) Assume $d_{tr}(G) = t$ and (D_1, D_2, \dots, D_t) is a partition of $V(G)$ into t tree dominating sets. Since each $\langle D_i \rangle$ is a tree dominating set, it follows that $|D_i| \geq \gamma_{tr}(G)$ for $i = 1, 2, \dots, t$. Hence, $n = \sum_{1 \leq i \leq t} |D_i| \geq t\gamma_{tr}(G)$. That is $d_{tr}(G) \leq \frac{n}{\gamma_{tr}(G)}$.

(3) Let S denote the cut set with cardinality $\kappa(G)$. Then $\langle V(G) - S \rangle$ is disconnected. It is obvious that any tree dominating set must contain at least one vertex of S . So G has at most $|S|$ pairwise disjoint tree dominating sets. Hence, $d_{tr}(G) \leq \kappa(G)$.

The sharpness of the bounds is obvious from the following two corollaries and Theorem 4. \square

Now, we give some exact values of the tree domatic number for some classes of graphs.

Corollary 1. (1) Let T denote a tree with $n \geq 3$. Then $d_{tr}(T) = 1$.

$$(2) d_{tr}(C_n) = \begin{cases} 3 & \text{for } n = 3, \\ 2 & \text{for } n = 4, \\ 0 & \text{for } n \geq 5. \end{cases}$$

$$(3) d_{tr}(K_{n,m}) = \min(n, m) \text{ for } n, m \geq 2.$$

Proof. (1) By Theorem 1, since $\kappa(T) = 1$, it follows that $d_{tr}(G) \leq 1$. It is obvious that T is a tree dominating set for T . So, $d_{tr}(T) = 1$.

(2) Let $C_n = v_1v_2 \dots v_nv_1$. Since $\gamma_{tr}(C_n) = n - 2$, from Theorem 1 it follows that $d_{tr}(T) \leq \frac{n}{n-2} = 1 + \frac{2}{n-2}$.

If $n = 3$, then $d_{tr}(T) \leq 3$. It is obvious that each vertex is a tree dominating set for C_3 . So, $d_{tr}(T) = 3$.

If $n = 4$, then $d_{tr}(T) \leq 2$. It is obvious that $\{v_1, v_2\}$ and $\{v_3, v_4\}$ are two pairwise disjoint tree dominating sets for C_4 . So, $d_{tr}(T) = 2$.

If $n \geq 5$, then $d_{tr}(T) \leq 1$. Hence, $V(C_n)$ is not a tree dominating sets for C_n . So, $d_{tr}(T) = 0$.

(3) Let $U = \{u_1, \dots, u_n\}$ and $W = \{w_1, \dots, w_m\}$ form a two vertex partition of $K_{n,m}$. Without loss of generality, assume $\min(n, m) = n$. By Theorem 1, it follows that $d_{tr}(K_{n,m}) \leq \delta(K_{n,m}) = n$. It is obvious that $\{u_i, w_i\}$ for $i = 1, \dots, n - 1$ and $\{u_n, v_n, \dots, v_m\}$ are n pairwise disjoint tree dominating sets of $K_{n,m}$. So, $d_{tr}(K_{n,m}) = n$. \square

Corollary 2. Let G be a connected graph with $\kappa(G) = 1$. If G contains a cycle, then $d_{tr}(G) = 0$.

Let H_t denote the set of trees of order t . Assume $T_1, T_2, \dots, T_k \in H_t$. Let $G(T_1, T_2, \dots, T_k)$ be obtained from T_1, T_2, \dots, T_k by joining each vertex v in T_i to exactly one vertex of each T_j for $j = 1, 2, \dots, k$ and $j \neq i$, and for arbitrary two vertices $u, v \in T_i$, $N(u) \cap V(T_j) \neq N(v) \cap V(T_j)$. Let $G_t = \{G(T_1, T_2, \dots, T_k) | T_1, T_2, \dots, T_k \in H_t\}$.

Theorem 2. For any $G(T_1, T_2, \dots, T_k) \in G_t$, there is $d_{tr}(G(T_1, T_2, \dots, T_k)) = k$.

Proof. It is clear that T_1, T_2, \dots, T_k is a tree domatic partition of $G(T_1, T_2, \dots, T_k)$. Hence, $d_{tr}(G(T_1, T_2, \dots, T_k)) \geq k$. Since $\delta(G(T_1, T_2, \dots, T_k)) = k$, it follows that $d_{tr}(G(T_1, T_2, \dots, T_k)) \leq k$ by Theorem 1. So, $d_{tr}(G(T_1, T_2, \dots, T_k)) = k$. \square

Theorem 3. Let G be a connected graph of order n with tree domatic number $k \geq 1$. Then G must have at least $\frac{(k+1)n}{2} - k$ edges. Furthermore, G has exactly $\frac{(k+1)n}{2} - k$ edges if and only if $n \equiv 0 \pmod{k}$ and $G \in G_{\frac{n}{k}}$.

Proof. Let D_1, D_2, \dots, D_k be a tree domatic partition of G and let $G_i = \langle D_i \rangle$. Without loss of generality, we assume that G_i has order n_i for $i = 1, 2, \dots, k$. Let $E(D_i, V(G) - D_i)$ denote such a set of edges that for any edge one of its endpoints belongs to D_i and the other belongs to $V(G) - D_i$. Since D_i is a tree dominating set for G , it follows that $|E(D_i, V(G) - D_i)| \geq n - n_i$. So,

$$\begin{aligned} |E(G)| &\geq \sum_{1 \leq i \leq k} |E(G_i)| + \frac{1}{2} \sum_{1 \leq i \leq k} |E(D_i, V(G) - D_i)| \geq \\ &\geq \sum_{1 \leq i \leq k} (n_i - 1) + \frac{1}{2} \sum_{1 \leq i \leq k} (n - n_i) \geq \\ &\geq (n - k) + \frac{1}{2}n(k - 1) = \\ &= \frac{(k + 1)n}{2} - k. \end{aligned}$$

For each $G \in G_{\frac{n}{k}}$, there exist T_1, T_2, \dots, T_k that belong to $H_{\frac{n}{k}}$ such that $G = G(T_1, T_2, \dots, T_k)$. By the definition of $G(T_1, T_2, \dots, T_k)$, it is clear that G has exactly $\frac{(k+1)n}{2} - k$ edges.

Conversely, if G has exactly $\frac{(k+1)n}{2} - k$ edges, then for each D_i , $|E(D_i, V(G) - D_i)| = n - n_i$. Since for each $v \in V(G) - D_i$, v is dominated by at least one vertex in D_i , it follows that v is dominated by exactly one vertex in D_i . Hence, $|D_1| = |D_2| = \dots = |D_k|$. Indeed, let us for the contrary assume that $|D_i| \neq |D_j|$. Without loss of generality, assume that $|D_i| > |D_j|$. Since each vertex $v \in D_i$ is dominated by exactly one vertex in D_j , it follows that there exists one vertex in D_j which is dominated by at least two vertices in D_i , which is a contradiction. So, $k|D_1| = n$. That is $n \equiv 0 \pmod{k}$. For arbitrary two vertices $u, v \in D_i$, it is clear that $N(u) \cap D_j \neq N(v) \cap D_j$. Let $T_i = \langle D_i \rangle$ for $i = 1, 2, \dots, k$. It is clear that $G = G(T_1, T_2, \dots, T_k)$. Hence, $G \in G_{\frac{n}{k}}$. \square

3. CHARACTERIZATION OF REGULAR GRAPHS AND PLANAR GRAPHS ATTAINING THE UPPER BOUNDS

Theorem 4. Let G be an r -regular graph of order n , where $r < n - 1$. Then $d_{tr}(G) = r$ if and only if $n = 2r$ and there exists a perfect matching $M = \{e_1, e_2, \dots, e_r\}$ of

G. If $e_i = u_i v_i$ for $i = 1, 2, \dots, r$, then for each u_i there is $|N(u_i) \cap e_j| = 1$ and $N(u_i) \cap e_j \neq N(v_i) \cap e_j$.

Proof. Let G be an r -regular graph with $d_{tr}(G) = r$ and let D_1, D_2, \dots, D_r be a tree domatic partition of G . By Theorem 3, it follows that $|E(G)| \geq \frac{(r+1)n}{2} - r$. That is $\frac{rn}{2} \geq \frac{(r+1)n}{2} - r$. So, $n \leq 2r$. If $n < 2r$, then there exists a D_i such that $|D_i| = 1$. Let $D_1 = \{v\}$. Since each vertex in D_j is dominated by v for $j = 2, \dots, r$, it follows that $d(v) = n - 1$ and G is a complete graph, which is a contradiction. Hence, $n = 2r$ and $|D_i| = 2$ for $i = 1, 2, \dots, r$. Let $M = \{\langle D_1 \rangle, \langle D_2 \rangle, \dots, \langle D_r \rangle\}$. Then M is a perfect matching of G . Assume $\langle D_i \rangle = e_i = u_i v_i$ for $i = 1, 2, \dots, r$. Since $d(u_i) = r$ and u_i is dominated by at least one vertex in D_j for $j = 1, 2, \dots, r$ and $j \neq i$, it follows that $|N(u_i) \cap e_j| = 1$. If $N(u_i) \cap e_j = N(v_i) \cap e_j$, that is both u_i and v_i are adjacent to the same vertex, say u_j , then $d_{\langle D_i \cup D_j \rangle}(u_j) = 3$. Hence, there exists a D_k such that u_j is not dominated by any vertex in D_k for $k \in \{1, 2, \dots, r\} \setminus \{i, j\}$, which is a contradiction.

Conversely, let $D_i = \{u_i, v_i\}$ for $i = 1, 2, \dots, r$. It is obvious that (D_1, D_2, \dots, D_r) is a tree domatic partition of G . Hence, $d_{tr}(G) \geq r$. By Theorem 1, it follows that $d_{tr}(G) = r$. \square

Corollary 3. *If G is a cubic graph and $d_{tr}(G) = 3$, then G is $K_{3,3}$ or $K_3 \times K_2$.*

Lemma 1. *A graph G is planar(outerplanar) if and only if neither K_5 nor $K_{3,3}$ (K_4 nor $K_{2,3}$) is a contraction of a subgraph of G .*

In a way similar to used in [3], we may prove the following theorem.

Theorem 5. *Let G be a planar graph of order n . Then the tree domatic number of G is at most 4 and K_4 is the only planar graph achieving this bound.*

Proof. Assume G is a planar graph of order n such that $d_{tr}(G) \geq 5$. By Theorem 3, it follows that $|E(G)| \geq 3n - 5$. But this contradicts the well known upper bound of $3n - 6$ for a planar graph of order n . Thus, $d_{tr}(G) \leq 4$.

Consider the case of $d_{tr}(G) = 4$ and G not isomorphic to K_4 . Let $D_1 \cup D_2 \cup D_3 \cup D_4$ be a tree domatic partition of G . Without loss of generality, assume $|D_1| \geq 2$. Let a and b be adjacent vertices in D_1 . Let H be the contraction of G formed by identifying the vertices in each of D_2, D_3 and D_4 to a single vertex x_2, x_3 and x_4 , respectively. H is planar since it is a contraction of a planar graph. But since D_2 is a dominating set in G , it follows that x_2 is adjacent to each of a, b, x_3, x_4 in H . Similar statements hold for x_3 and x_4 . But then H contains the subgraph $\langle \{a, b, x_2, x_3, x_4\} \rangle$ which has been shown to be isomorphic to K_5 , which is a contradiction. \square

Theorem 6. *Let G be a planar graph such that $d_{tr}(G) = 3$ and let $D_1 \cup D_2 \cup D_3$ be any tree domatic partition of G . Then each of the induced subgraphs $\langle D_1 \rangle, \langle D_2 \rangle$ and $\langle D_3 \rangle$ is a path.*

Proof. Let G be a planar graph and $D_1 \cup D_2 \cup D_3$ be any tree domatic partition of G such that $\langle D_1 \rangle$ has a vertex a of degree at least three. Let b, c, d be three of its neighbours in D_1 and let $u_2 \in D_2$ and $u_3 \in D_3$. Let H be the planar graph obtained

from G by contracting D_i into u_i for $i = 2, 3$, and by removing any of the edges bc, bd, cd which is present in G . Since each vertex in D_1 is dominated by at least one vertex in D_2 , each of a, b, c and d is adjacent to u_2 in H . Consider any planar embedding of H . By Euler's formula, the induced subgraph $K = \langle \{a, b, c, d, u_2\} \rangle$ has four regions, and u_2 belongs to the boundary of each of these regions. Since, K contains a subgraph isomorphic to $K_{2,3}$, it is not outerplanar, and so none of the four regions has a boundary which contains all of a, b, c and d . Now u_3 lies in one of these regions and can be adjacent only to vertices on the boundary of this region. But then there exists in $\{a, b, c, d\}$ a vertex which it is not dominated by u_3 , contradicting our assumption above. Therefore, the maximum degree of any of the induced subgraphs $\langle D_1 \rangle$, $\langle D_2 \rangle$ and $\langle D_3 \rangle$ is no more than two. Hence, each of the induced subgraphs $\langle D_1 \rangle$, $\langle D_2 \rangle$ and $\langle D_3 \rangle$ is a path. \square

Theorem 7. *Suppose $2 \leq r \leq s$. If $r = 2$, then $d_{tr}(G_{r,s}) = 2$; otherwise, $d_{tr}(G_{r,s}) = 0$.*

Proof. Let P_r be the path v_1, v_2, \dots, v_r and let P_s be the path w_1, w_2, \dots, w_s .

The case $r = 2$ is easy. Assume $r \geq 3$ and $\{D_1, D_2\}$ is a partition of $V(G_{r,s})$ into two tree dominating sets. For ease of reference, let a, b, c, d, e, f be the vertices (v_1, w_2) , (v_2, w_1) , (v_1, w_{s-1}) , (v_2, w_s) , (v_{r-1}, w_s) , (v_r, w_{s-1}) , respectively. Note that if $r = 3$, then $d = e$. (If, in addition, $s = 3$, then $a = c$ as well). We may assume without loss of generality that $a \in D_1$ and $b \in D_2$. Since D_1 and D_2 are tree dominating sets of $G_{r,s}$, some neighbours of (v_1, w_s) must belong to D_2 , and some neighbours of (v_r, w_s) must belong to D_1 . Let $\{x\} = D_2 \cap \{c, d\}$ and let $\{y\} = D_1 \cap \{e, f\}$.

If $r = 3 = s$, then $f \in D_1$ and $d \in D_2$. But then $(v_2, w_2) \in D_1$ and so $\langle D_2 \rangle$ is not connected, a contradiction. Therefore, assume $s \geq 4$. There are three cases to consider. If $r = 3$ and $y = e$, then it follows that $x = c$ and $f \in D_2$. Thus the vertices of any $a - y$ path form a cutset which separates vertices b and c . If $r = 3$ and $y = f$, then the vertex set of any $a - y$ path separates vertices b and d . If $r \geq 4$, then vertices b and x are separated by the vertex set of any $a - y$ path. In each case, we come to a contradiction with the assumption that $\langle D_2 \rangle$ is connected. So, $d_{tr}(G_{r,s}) \leq 1$. Since $G_{r,s}$ is not a tree, it follows that $d_{tr}(G_{r,s}) = 0$. \square

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