

Performance Evaluation of Two-Link Multirate Loss Models with Restricted Accessibility

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Abstract—We consider a two-link communication system with restricted accessibility that services Poisson arriving calls of many service-classes and propose a multirate teletraffic loss model for its analysis. In a restricted accessibility system, call blocking occurs even if available resources do exist at the time of a call's arrival. In the two-link system under consideration, each link has two thresholds (offloading and support) which express the in-service calls in a link. The offloading threshold represents the point from which a link offloads calls. The support threshold (which is lower than the offloading threshold) defines the point up to which a link supports offloaded calls. The two-link system with restricted accessibility is modeled as a loss system whose steady state probabilities do not have a product form solution. However, approximate formulas for the determination of call blocking probabilities are proposed. In addition, we also provide a corresponding analysis related to the case of quasi-random traffic (i.e. traffic generated by a finite number of users). The accuracy of all formulas is verified through simulation and is found to be quite satisfactory.

Keywords—accessibility, blocking, non-product form, Poisson, quasi-random, threshold.

1. Introduction

Bandwidth sharing policies are quality of service (QoS) guarantee mechanisms that are necessary for the provision of bandwidth required by calls in a communication link. Assuming that the link is modeled as a loss system carrying call-level traffic, the most common bandwidth sharing policy is the complete sharing (CS) policy. In the CS policy, a new call is blocked if its required bandwidth units (b.u.) exceed the link's available b.u. In addition to the "CS policy" term adopted in this paper, other equivalent terms are also used in the literature, such as "full accessibility" or "full availability" [1], [2]. The latter, however, usually refers to the proportion of time over which the link is available [3].

The simplest loss system that adopts the CS policy is the Erlang loss system [4]. In this system, with its analysis based

on the Erlang loss model, new calls follow a Poisson process, require one b.u. in order to be accepted by the system and have a generally distributed service time. Call blocking occurs if all b.u. are occupied at the time arrival of a given call. The fact that call blocking probabilities (CBP) are determined according to the Erlang B formula has led to an extensive amount of Erlang loss model extensions for the call-level analysis of wired (e.g. [5]–[20]), wireless (e.g. [21]–[33]), satellite (e.g. [34]–[36]) and optical networks (e.g. [37]–[43]).

In [29], a two-link loss system servicing Poisson traffic is considered and studied. Arriving calls belong to a single service-class and each call requests a single b.u. in order to be accepted in a link. Each link may service calls offloaded from the other link. An offloaded call is a new call that arrives in a link but will be served by the other link, subject to bandwidth availability. This offloading mechanism operates with the aid of a high and a low threshold per link, expressing the number of calls serviced by each link. The high threshold is the offloading threshold, while the low threshold is the support threshold (see Fig. 1). The latter expresses the point up to which the link is capable of supporting offloaded calls (from the other link). The offloading threshold determines the point from which call offloading between the two links may start.

Due to the offloading mechanism, there is no local balance (LB) between adjacent states and, therefore, the steady state probabilities of this system do not have a product form solution (PFS) (see the tutorial example of [44]). Thus, the CBP determination can be based either on the accurate but complex method of solving a set of linear global balance (GB) equations, or on an approximate but efficient method that relies on the Erlang B formula and on the assumption (approximation) that the links are independent.

Such an offloading scheme may find a potential application in mobile/Wi-Fi networks. To manage the increasing traffic in mobile networks, traffic may be offloaded to Wi-Fi networks [45], [46]. In order to increase the bandwidth of Wi-Fi access links, recent research focuses on bandwidth

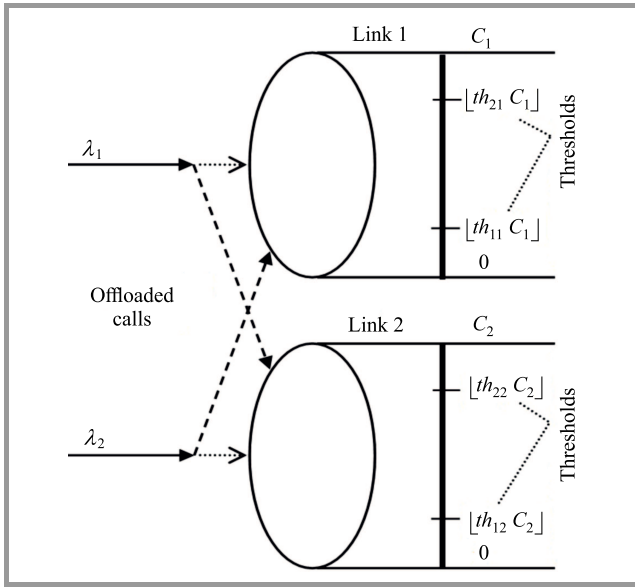


Fig. 1. The two-link loss system.

sharing policies that should be adopted and on the aggregation of backhaul access link capacities. The impact that such aggregation exerts on CBP may be studied via the single-rate model of [29].

In this paper, we extend the model of [29] by considering that the system accommodates Poisson arriving calls of numerous service-classes. Nowadays, this consideration is a sine-qua-non condition in multidimensional network traffic environments. In addition, in the proposed new model we incorporate the notion of restricted accessibility – not only in the case of Poisson traffic, but also in the case of quasi-random traffic (traffic generated by a finite number of users). In a restricted accessibility system, call blocking may occur even if b.u. are available at the time of arrival of a given call. The term “restricted accessibility” covers the following:

- Bandwidth sharing policies, such as the bandwidth (trunk) reservation policy [5], [10], [14], the threshold policy [16], [18] or the probabilistic threshold policy [30], [47]. In the bandwidth reservation policy, call blocking can occur if the available b.u. of the system are reserved at the time of an arrival of a call. In the threshold and the probabilistic threshold policies, a predefined threshold (different for each service-class) is set in order to express the number of in-service calls (of each service-class). If the acceptance of a new call leads to a value that is above that threshold, then call blocking always occurs (threshold policy) or it occurs with a certain probability (probabilistic threshold policy).
- The case where each state of the system (excluding the state where there are no calls in the system) is associated with a blocking probability. Such an approach may be useful when modeling interference between neighboring cells (e.g. in CDMA sys-

tems) [3], [48]. In this paper, we focus on this type of restricted accessibility and propose an approximate method for the CBP calculation which is verified via simulation and is found to be quite satisfactory. The CBP calculation in the proposed two-link model is based on the Erlang multirate loss model (EMLM) [49], [50] which refers to a link that services Poisson traffic generated from different service-classes.

In the remainder of this paper, in Section 2, we review the model of [29]. In Section 3, we propose the extension of [29] which includes the case of many service-classes, as well as the notion of restricted accessibility. In Section 4, we present the corresponding analytical model for the case of multirate quasi-random traffic. In Section 5, we provide analytical and simulated CBP results for the proposed model, assuming the existence of Poisson traffic. We conclude in Section 6.

2. Review of the Two-link Single-rate System

We consider a two-link system of capacities C_1 and C_2 b.u. Each link services single-rate Poisson traffic. Arriving calls require one b.u. in order to be accepted in the system and have a generally distributed service-time with a mean of μ^{-1} . Let λ_l be the call arrival rate in link l ($l = 1, 2$) and let j_l be the occupied b.u. in link l . Then, $0 \leq j_1 \leq C_1$ and $0 \leq j_2 \leq C_2$. Note that j_l expresses the number of in-service calls in link l , since each call requires one b.u.

Each link l ($l = 1, 2$) has a support (low) threshold th_{1l} and an offloading (high) threshold th_{2l} , with $th_{1l} < th_{2l}$ and $0 \leq th_{1l}, th_{2l} \leq 1$. By denoting the largest integer not exceeding x as $\lfloor x \rfloor$ and based on Fig. 1, the role of these thresholds in link l is described in the following manner:

- If $0 \leq j_l < \lfloor th_{1l} C_l \rfloor$, then link l is in a *support mode* of operation. In that mode, the link can service new calls that arrive in link l and offloaded calls from link m ($m = 1, 2, m \neq l$).
- If $\lfloor th_{1l} C_l \rfloor \leq j_l < \lfloor th_{2l} C_l \rfloor$, then link l is in a *normal mode* of operation. In that mode, the link does not service offloaded calls from the other link.
- If $\lfloor th_{2l} C_l \rfloor \leq j_l$, then link l operates in an *offloading mode*. In that mode, a new call that initially arrives in link l is offloaded to link m . If that link is in *support mode* (i.e. $0 \leq j_m < \lfloor th_{1m} C_m \rfloor$), then the call is accepted in link m . Otherwise, if $j_l \leq C_l - 1$, the call is accepted in link l , whereas if $j_l > C_l - 1$, the call is blocked and lost.

Based on the above description, the call admission mechanism applicable to a call that arrives in link l ($l = 1, 2$) consists of two steps:

1. If $(0 \leq j_l < \lfloor th_{2l}C_l \rfloor)$, then the call is serviced via link l .
2. If $\lfloor th_{2l}C_l \rfloor \leq j_l$, then:
 - If $0 \leq j_m < \lfloor th_{1m}C_m \rfloor$, the call is offloaded to link m .
 - If $\lfloor th_{1m}C_m \rfloor \leq j_m$, then link m does not support offloaded calls from link l since it operates in *normal mode*. In that case, the call will be handled by link l . Thus, if $j_l \leq C_l - 1$, the call is accepted in link l . Otherwise, call blocking occurs.

Due to the *offloading* and *support modes* modes of the two links, there is no LB between adjacent system states and, therefore, the steady state distribution, $P(j) = P(j_1, j_2)$, of such a system cannot be described by a PFS. To determine $P(j_1, j_2)$, two methods exist in the literature.

The first method provides accurate CBP results (compared to simulation results) but is quite complex, since it requires the solution of a set of linear GB equations for each state $\mathbf{j} = (j_1, j_2)$ expressed as *rate into state \mathbf{j} = rate out of state \mathbf{j}* :

$$\begin{aligned} & \lambda_1(j_1 - 1, j_2)P(j_1 - 1, j_2) + \lambda_2(j_1, j_2 - 1)P(j_1, j_2 - 1) \\ & + (j_1 + 1)\mu P(j_1 + 1, j_2) + (j_2 + 1)\mu P(j_1, j_2 + 1) \\ & = \lambda_1(j_1, j_2)P(j_1, j_2) + \lambda_2(j_1, j_2)P(j_1, j_2) \\ & \quad + (j_1\mu + j_2\mu)P(j_1, j_2), \quad (1) \end{aligned}$$

where:

$$l = 1, 2, \quad m \neq l$$

and

$$\lambda_l(j_1, j_2) = \begin{cases} \lambda_l + \lambda_m & \text{if } (j_l < \lfloor th_{1l}C_l \rfloor) \cap (j_m \geq \lfloor th_{2m}C_m \rfloor) \\ 0 & \text{if } (j_l \geq \lfloor th_{2l}C_l \rfloor) \cap (j_m < \lfloor th_{1m}C_m \rfloor) \\ 0 & \text{if } (j_1, j_2) \text{ is a boundary state} \\ \lambda_l & \text{otherwise} \end{cases}. \quad (2)$$

Having obtained $P(j_1, j_2)$, we can determine CBP in each link, P'_{b_1} and P'_{b_2} via Eqs. (3) and (4), respectively [29]:

$$P'_{b_1} = \sum_{j_2=\lfloor th_{12}C_2 \rfloor}^{C_2} P(C_1, j_2), \quad (3)$$

$$P'_{b_2} = \sum_{j_1=\lfloor th_{11}C_1 \rfloor}^{C_1} P(j_1, C_2). \quad (4)$$

Equation (3) expresses the fact that call blocking occurs in the first link if all b.u. are occupied (i.e. if $j_1 = C_1$) and, at the same time, the other link does not operate in the support mode (i.e. if $\lfloor th_{12}C_2 \rfloor \leq j_2$). Equation (4) may be interpreted accordingly.

To determine CBP in the system of [29], the following weighted summation can be used:

$$P'_b = \frac{\lambda_1}{\lambda_1 + \lambda_2} P'_{b_1} + \frac{\lambda_2}{\lambda_1 + \lambda_2} P'_{b_2}. \quad (5)$$

Contrary to the first method, the second method is simpler but provides approximate CBP results. The approximation lies on the fact that each link l is modeled as an independent Erlang loss system of capacity C_l ($l = 1, 2$).

The CBP in each link may be approximated by Eqs. (6) and (7), respectively:

$$P_{b_1} = P_1(C_1) P_2(j_2 \geq \lfloor th_{12}C_2 \rfloor), \quad (6)$$

$$P_{b_2} = P_2(C_2) P_1(j_1 \geq \lfloor th_{11}C_1 \rfloor), \quad (7)$$

where $P_l(C_l)$ is the CBP in link l ($l = 1, 2$).

The values of $P_l(C_l)$ in Eqs. (6) and (7) may be determined via the Erlang B formula (see Eq. (8a) for the closed form or Eq. (8b) for the recurrent form):

$$P_l(C_l) = \frac{\frac{a_l^{C_l}}{C_l!}}{\sum_{i=0}^{C_l} \frac{a_l^i}{i!}}, \quad a_l = \frac{\lambda_l}{\mu}, \quad (8a)$$

$$P_l(C_l) = \frac{a_l P_l(C_l - 1)}{C_l + a_l P_l(C_l - 1)}, \quad C_l \geq 1, \quad P_l(0) = 1. \quad (8b)$$

As far as the values of $P_l(j_l \geq \lfloor th_{1l}C_l \rfloor)$ in Eqs. (6) and (7) are concerned, they can be determined by:

$$P_l(j_l \geq \lfloor th_{1l}C_l \rfloor) = \sum_{j_l=\lfloor th_{1l}C_l \rfloor}^{C_l} P_l(j_l), \quad (9)$$

where $P_l(j_l)$ is calculated according to the truncated Poisson distribution:

$$P_l(j_l) = \frac{\frac{a_l^{j_l}}{j_l!}}{\sum_{i=0}^{C_l} \frac{a_l^i}{i!}}, \quad a_l = \frac{\lambda_l}{\mu}. \quad (10)$$

As far as the total blocking probability in the two-link system is concerned, it can be determined via Eq. (5), where P'_{b_1} and P'_{b_2} are replaced by P_{b_1} and P_{b_2} , determined in Eqs. (6) and (7), respectively.

An additional recursive way for the determination of $P_l(j_l)$, $j_l = 1, \dots, C_l$ is based on the link independence assumption. In the Erlang loss model, used to describe each link l , there exist the following LB between states $j_l - 1$ and j_l [44]:

$$j_l P'_l(j_l) = a_l P'_l(j_l - 1). \quad (11)$$

Based on Eq. (11), we can determine the unnormalized values of $P'_l(j_l)$'s considering an initial value of $P'_l(0) = 1$. Then, the normalized values of $P'_l(j_l)$'s are given by:

$$P_l(j_l) = \frac{P'_l(j_l)}{\sum_{x=0}^{C_l} P'_l(x)}. \quad (12)$$

Based on Eq. (12), we can compute P_{b_1} , P_{b_2} and, consequently, the total CBP, via Eqs. (6), (7) and (5), respectively.

3. The Proposed Multirate Loss Model – Poisson Case

In the proposed loss model, each link services Poisson arriving calls of K service-classes. New calls of service class k ($k = 1, \dots, K$) require b_k b.u. in order to be accepted in a link and have a generally distributed service-time with a mean of μ_k^{-1} . Let λ_{1k} and λ_{2k} be the arrival rates in each link of service-class k calls, respectively. We also denote, by j_1 and j_2 , the b.u. occupied in each link. Similarly to Section 2, each link l ($l = 1, 2$) has a support threshold th_{1l} and an offloading threshold th_{2l} , with $th_{1l} < th_{2l}$ and $0 \leq th_{1l}, th_{2l} \leq 1$.

To incorporate restricted accessibility into our model, we assume that each state j_l of link l , except for state $j_l = 0$ where link l is empty, is associated with a blocking probability, $pb_{l,k}(j_l)$. When there are no available b.u. for calls of service-class k in link l (i.e. when $j_l \geq C_l - b_k + 1$), then $pb_{l,k}(j_l) = 1$. Similarly, in the case of an empty system, $pb_{l,k}(0) = 0$.

The procedure of admitting a new service class k call that arrives in link l ($l = 1, 2$) is the following:

1. If $(0 \leq j_l < \lfloor th_{2l}C_l \rfloor)$, then the call is handled by link l . In addition, if $j_l + b_k \leq C_l$, then the call is accepted in link l with probability $1 - pb_{l,k}(j_l)$.
2. If $\lfloor th_{2l}C_l \rfloor \leq j_l$, then:
 - If $(0 \leq j_m < \lfloor th_{1m}C_m \rfloor)$, the call is offloaded to link m and if $j_m + b_k \leq C_m$, the call is accepted in link m with probability $1 - pb_{m,k}(j_m)$.
 - If $\lfloor th_{1m}C_m \rfloor \leq j_m$, then link m operates in *normal mode* and, therefore, does not support offloaded calls. Thus, the call is handled by link l . If $j_l + b_k \leq C_l$, then the call is accepted in link l with probability $1 - pb_{l,k}(j_l)$. Otherwise, the call is blocked and lost.

To calculate the CBP of service-class k calls, we assume that each link is an independent EMLM system under restricted accessibility [48] and, therefore, the CBP of service-class k calls in the first and the second link can be given via Eqs. (13) and (14), respectively:

$$P_{\text{res},b_{1k}} = P_{\text{res},1k}(C_1)P_{\text{res},2}(j_2 \geq \lfloor th_{12}C_2 \rfloor), \quad (13)$$

$$P_{\text{res},b_{2k}} = P_{\text{res},2k}(C_2)P_{\text{res},1}(j_1 \geq \lfloor th_{11}C_1 \rfloor), \quad (14)$$

where $P_{\text{res},lk}(C_{lk})$ is the CBP of service-class k calls in link l ($l = 1, 2$) and $P_{\text{res},l}(j_l \geq \lfloor th_{1l}C_l \rfloor)$ is the probability that link l does not operate in the support mode.

The values of $P_{\text{res},lk}(C_l)$ in Eqs. (13) and (14) can be given by:

$$P_{\text{res},lk}(C_l) = \sum_{j_l=1}^{C_l} G_l^{-1} q(j_l) pb_{l,k}(j_l), \quad (15)$$

where $q(j_l)$ expresses the unnormalized values of the occupancy distribution of link l ($l = 1, 2$) while $G_l = \sum_{j_l=0}^{C_l} q(j_l)$ refers to the normalization constant.

In Eq. (15), the values of $q(j_l)$ can be computed via:

$$q(j_l) = \begin{cases} 1 & \text{for } j_l = 0 \\ \frac{1}{j_l} \sum_{k=1}^K a_{lk} b_k q(j_l - b_k) \times \\ \quad [1 - pb_{l,k}(j_l - b_k)] & \text{for } j_l = 1, \dots, C_l \end{cases}, \quad (16)$$

where $a_{lk} = \lambda_{lk}/\mu_k$ is the total offered traffic-load of service-class k calls in link l .

Regarding the values of $P_{\text{res},l}(j_l \geq \lfloor th_{1l}C_l \rfloor)$, in (13) and (14), they can be calculated by:

$$P_{\text{res},l}(j_l \geq \lfloor th_{1l}C_l \rfloor) = \sum_{j_l=\lfloor th_{1l}C_l \rfloor}^{C_l} G_l^{-1} q(j_l), \quad (17)$$

where $q(j_l)$ is given by (16).

Finally, Eq. (18) is proposed for determining the total blocking probability of service-class k calls in the two-link system:

$$P_{\text{res},b_k} = \frac{\lambda_{1k}}{\lambda_{1k} + \lambda_{2k}} P_{\text{res},b_{1k}} + \frac{\lambda_{2k}}{\lambda_{1k} + \lambda_{2k}} P_{\text{res},b_{2k}}. \quad (18)$$

4. The Proposed Multirate Loss Model – Quasi-Random Case

Contrary to the model proposed in Section 3, we now assume that each link services quasi-random traffic generated by K service-classes. New calls of service-class k ($k = 1, \dots, K$) in link l ($l = 1, 2$) are generated via a finite source population N_{lk} and require b_k b.u. Let $\lambda_{1k,\text{fin}}$ and $\lambda_{2k,\text{fin}}$ be the arrival rates of service-class k idle sources in the first and second link, respectively. Then, $\lambda_{lk,\text{fin}} = (N_{lk} - n_{1k} - n_{2k})v_{lk}$, where n_{lk} refers to the in-service calls of service-class k in link l and v_{lk} is the arrival rate per idle source of service-class k in link l . The corresponding offered traffic-load per idle source is $a_{lk,\text{idle}} = v_{lk}/\mu_k$. Note that if $N_{lk} \rightarrow \infty$ for all service-classes and the total offered traffic-load is constant, then calls arrive in the system according to a Poisson process, resulting in the model described in Section 3.

The process of admitting a new service-class k call that arrives in link l ($l = 1, 2$) is similar to that described in Section 3 and is therefore omitted.

To calculate the time congestion probabilities of service-class k calls, we assume that each link is an independent Engset multirate loss model (EnMLM) under restricted accessibility and, therefore, time congestion probabilities of service-class k calls in the first and the second link can be determined via Eqs. (19) and (20), respectively:

$$P_{\text{fin-res},b_{1k}} = P_{\text{fin-res},1k}(C_1)P_{\text{fin-res},2}(j_2 \geq \lfloor th_{12}C_2 \rfloor), \quad (19)$$

$$P_{\text{fin-res},b_{2k}} = P_{\text{fin-res},2k}(C_2)P_{\text{fin-res},1}(j_1 \geq \lfloor th_{11}C_1 \rfloor), \quad (20)$$

where $P_{\text{fin-res},lk}(C_{lk})$ is the time congestion probability of service-class k calls in link l ($l = 1, 2$) and

$P_{\text{fin-res},l}(j_l \geq \lfloor th_{1l}C_l \rfloor)$ is the probability that link l is not in the support mode.

The values of $P_{\text{fin-res},lk}(C_l)$ in Eqs. (19) and (20) can be determined via the following formula:

$$P_{\text{fin-res},lk}(C_l) = \sum_{j_l=1}^{C_l} G_l^{-1} q_{\text{fin}}(j_l) p b_{l,k}(j_l), \quad (21)$$

where $q_{\text{fin}}(j_l)$ expresses the unnormalized values of the occupancy distribution of link l ($l = 1, 2$), while $G_l = \sum_{j_l=0}^{C_l} q_{\text{fin}}(j_l)$ is the corresponding normalization constant. In Eq. (21), the values of $q_{\text{fin}}(j_l)$ can be calculated as follows:

$$q_{\text{fin}}(j_l) = \begin{cases} 1 & \text{for } j_l = 0 \\ \frac{1}{j_l} \sum_{k=1}^K (N_{lk} - Y_k) a_{lk, \text{idle}} b_k q_{\text{fin}} \times \\ \quad \left((j_l - b_k) [1 - p b_{l,k}(j_l - b_k)] \right) & \text{for } j_l = 1, \dots, C_l \end{cases}, \quad (22)$$

where $Y_k = y_{1k}(j_1 - b_k) - y_{2k}(j_2 - b_k)$ and $y_{lk}(j_l)$ is the average number of service-class k calls in state j_l of link l , assuming that the system accommodates Poisson traffic.

The values of $y_{lk}(j_l)$ are given by:

$$y_{lk}(j_l) = \frac{a_{lk} q(j_l - b_k) [1 - p b_{l,k}(j_l - b_k)]}{q(j)}, \quad (23)$$

where the values of $q(j)$ are computed via Eq. (16) (i.e. by the corresponding Poisson model).

The rationale behind Eqs. (22) and (23) is similar to that of the model from [51] that proposes an algorithm for the approximate determination of time congestion probabilities in the EnMLM.

Finally, regarding the values of $P_{\text{fin-res},l}(j_l \geq \lfloor th_{1l}C_l \rfloor)$, in Eqs. (19) and (20), they are given by:

$$P_{\text{fin-res},l}(j_l \geq \lfloor th_{1l}C_l \rfloor) = \sum_{j_l=\lfloor th_{1l}C_l \rfloor}^{C_l} G_l^{-1} q_{\text{fin}}(j_l), \quad (24)$$

where $q_{\text{fin}}(j_l)$ is determined via Eq. (22).

5. Numerical Examples

In this section, we consider two examples and provide simulation and analytical CBP results of the proposed model assuming Poisson traffic. Simulation results are based on Simscript III [52] and are mean values of 7 runs. In each run, ten million calls are generated. The first 5% of these generated calls are not considered in the CBP results so as to account for a warm-up period.

In the first example, we consider a system with the capacities of $C_1 = 24$ b.u. and $C_2 = 20$ b.u., accommodating two service-classes whose calls require $b_1 = 1$ and $b_2 = 2$ b.u., respectively. Let $\lambda_{11} = 9$ calls/min and $\lambda_{12} = 1$ call/min, for the first link. Similarly, let $\lambda_{21} = 7$ calls/min and

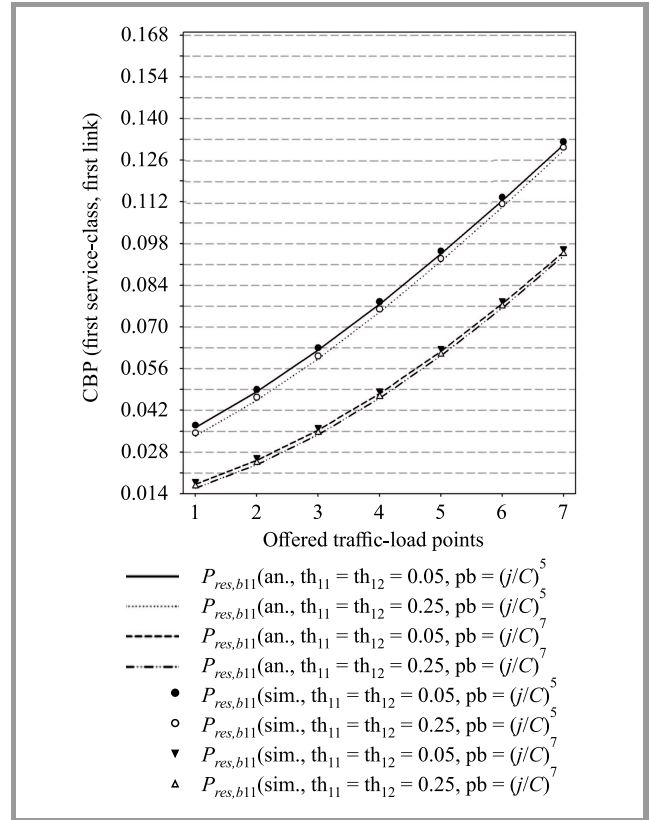


Fig. 2. CBP of the first service-class in the first link (example 1).

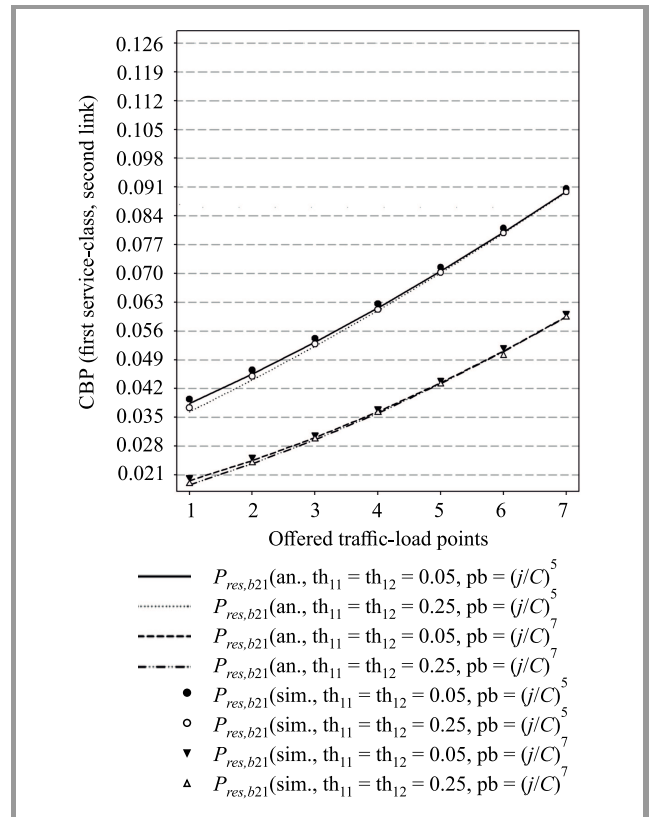


Fig. 3. CBP of the first service-class in the second link (example 1).

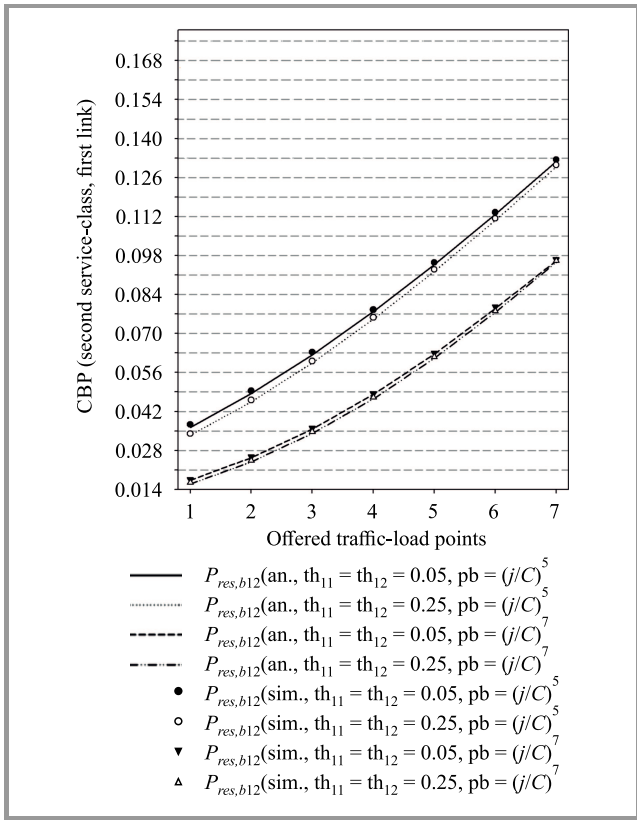


Fig. 4. CBP of the second service-class in the first link (example 1).

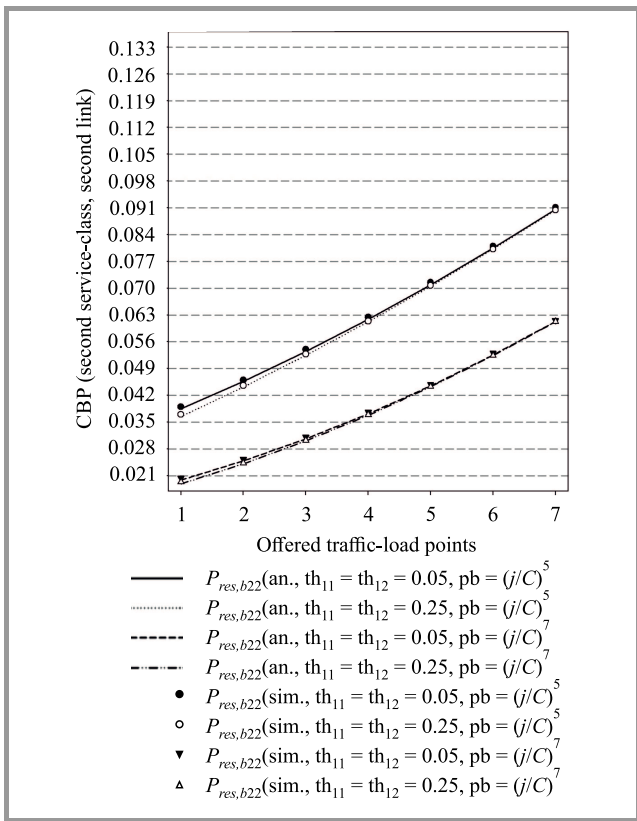


Fig. 5. CBP of the second service-class in the second link (example 1).

$\lambda_{22} = 1$ call/min, for the second link. Also let $\mu_1^{-1} = \mu_2^{-1} = 1.0$ min. Regarding the values of the thresholds, let the offloading thresholds equal $th_{21} = th_{22} = 0.7$, and let us consider two support threshold scenarios: (i) $th_{11} = th_{12} = 0.05$ and (ii) $th_{11} = th_{12} = 0.25$. Finally, regarding the restricted accessibility factors for each link, two sets are studied: (i) $pb_{l,1}(j_l) = pb_{l,2}(j_l) = (j_l/C_l)^5$ and (ii) $pb_{l,1}(j_l) = pb_{l,2}(j_l) = (j_l/C_l)^7$ where $l = 1, 2$.

In the x -axis of Figs. 2–5, λ_{11} and λ_{21} increase in steps of 1.0 and 0.5, respectively. So, point 1 is: $(\lambda_{11} = 9.0, \lambda_{12} = 1.0, \lambda_{21} = 7.0, \lambda_{22} = 1.0)$, while point 7 is: $(\lambda_{11} = 15.0, \lambda_{12} = 1.0, \lambda_{21} = 10.0, \lambda_{22} = 1.0)$.

In Figs. 2–3, we present the CBP for the first service-class in each link, respectively. In Figs. 4–5, the corresponding CBP results for the second service-class are presented. Figures 2–5 show that the analytical CBP results:

- Are close to the simulation results, especially when the values of support thresholds th_{11} and th_{12} are at a reasonable level (e.g. 0.05 to 0.25). Depending on the system, higher values of th_{11} and th_{12} may increase the discrepancy between simulation and analytical CBP results. This behavior may also be observed in [29] and is anticipated due to the fact that both links work independently.
- The choice of $pb_{l,k}(j_l)$ greatly affects CBP. The higher values of set 1 leads to much higher CBP compared to the values of set 2.

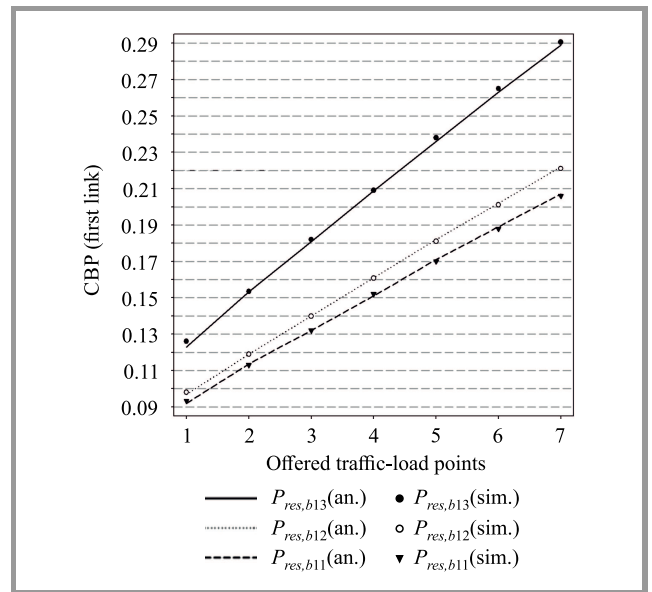


Fig. 6. CBP of each service-class in the first link (example 2).

In the second example, a larger system is considered. More precisely, we study a system of capacities $C_1 = 40$ b.u. and $C_2 = 45$ b.u. that services three service-classes whose calls require $b_1 = 1, b_2 = 3$ and $b_3 = 6$ b.u., respectively. Let $\lambda_{11} = 5$ calls/min, $\lambda_{12} = 3$ calls/min and $\lambda_{13} = 2$ calls/min, for the first link. Similarly, let $\lambda_{21} = 5$ calls/min, $\lambda_{22} = 3$ calls/min and $\lambda_{23} = 2$ calls/min, for the second link. Also

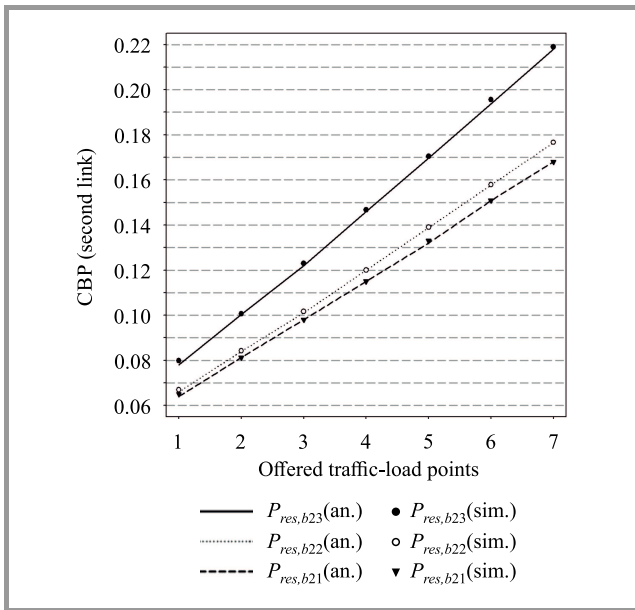


Fig. 7. CBP of each service-class in the second link (example 2).

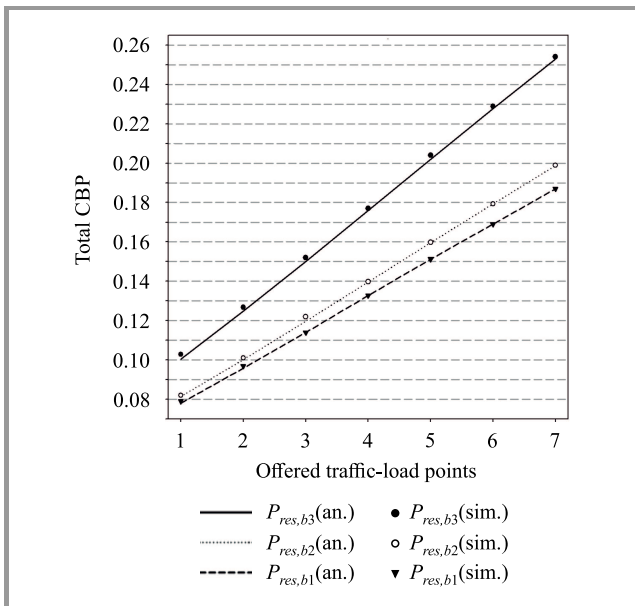


Fig. 8. Total CBP (example 2).

let $\mu_1^{-1} = \mu_2^{-1} = \mu_3^{-1} = 1$ min. Regarding the values of the thresholds, let the offloading thresholds equal $th_{21} = th_{22} = 0.7$ and let the support thresholds equal $th_{11} = th_{12} = 0.2$. Finally, regarding the restricted accessibility factors, let $pb_{l,1}(j_l) = pb_{l,2}(j_l) = (j_l/C_l)^7$ where $l = 1, 2$.

In the x -axis of Figs. 6–8, $\lambda_{11}, \lambda_{12}, \lambda_{13}, \lambda_{21}, \lambda_{22}$ and λ_{23} increase in steps of 0.2, respectively. So, point 1 is: $(\lambda_{11} = 5.0, \lambda_{12} = 3.0, \lambda_{13} = 2.0, \lambda_{21} = 5.0, \lambda_{22} = 3.0, \lambda_{23} = 2.0)$, while point 7 is: $(\lambda_{11} = 6.2, \lambda_{12} = 4.2, \lambda_{13} = 3.2, \lambda_{21} = 6.2, \lambda_{22} = 4.2, \lambda_{23} = 3.2)$. In Figs. 6–7, we show the CBP for all service-classes in each link, respectively. In Fig. 8, we present the total CBP results for all service-classes. Figures 6–8 show that the analytical CBP results are again quite close to the corresponding simulation re-

sults. A similar degree of accuracy has been observed for various two-link systems that we studied.

6. Conclusion

We propose new multirate loss models for the call-level analysis of a two-link system with restricted accessibility that accommodates Poisson or quasi-random arriving calls of different service-classes. In this system, each link may support calls offloaded from the other link. The proposed models do not have a PFS for the steady state probabilities due to the restricted accessibility and existence of the offloading mechanism. However, we show that an approximate method does exist for the determination of blocking probabilities, achieving a satisfactory degree of accuracy compared to simulation. As a future work, we intend to analyze the case of interference between the two links, using the proposed model as a springboard for further considerations.

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